

# Mastodon Theorem - 20 Years in the Making

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Combinatorics and Geometry Days III - MIPT, December 4, 2020

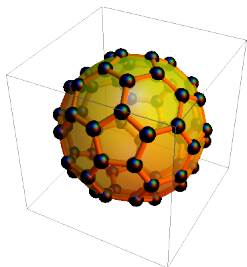


\*Jointly with Oleg Musin

# Optimal $s$ -energy and Log-optimal codes

**Thomson Problem (1904)** -  
 (“plum pudding” model of an atom)

Find the (most) stable (ground state) energy configuration (**code**) of  $N$  classical electrons (Coulomb law) constrained to move on the sphere  $\mathbb{S}^2$ .



**Generalized Thomson Problem ( $1/r^s$  potentials and  $\log(1/r)$ )**

A code  $C := \{\mathbf{x}_1, \dots, \mathbf{x}_N\} \subset \mathbb{S}^2$  that minimizes **Riesz  $s$ -energy**

$$E_s(C) := \sum_{j \neq k} \frac{1}{|\mathbf{x}_j - \mathbf{x}_k|^s}, \quad s > 0, \quad E_{\log}(w_N) := \sum_{j \neq k} \log \frac{1}{|\mathbf{x}_j - \mathbf{x}_k|}$$

is called an **optimal  $s$ -energy code** (**log-optimal** for  $s = 0$ )

# Optimal $s$ -energy codes on $\mathbb{S}^2$

## Known optimal $s$ -energy codes on $\mathbb{S}^2$

- $s = \log$ , Whyte's problem (1952, Monthly) ( $N = 2 - 6, 12$ );
- $s = 1$ , Thomson Problem (known for  $N = 2 - 6, 12$ )
- $s = -1$ , Fejes-Toth Problem (known for  $N = 2 - 6, 12$ )
- $s \rightarrow \infty$ , Tammes Problem (known for  $N = 1 - 12, 13, 14, 24$ )

## Limiting case - Best packing

For fixed  $N$ , any limit as  $s \rightarrow \infty$  of optimal  $s$ -energy codes is an optimal (maximal) code.

## Universally optimal codes

The codes with cardinality  $N = 2, 3, 4, 6, 12$  are special (*sharp codes*) and minimize large class of potential energies. First "non-sharp" is  $N = 5$  and very little is rigorously proven.

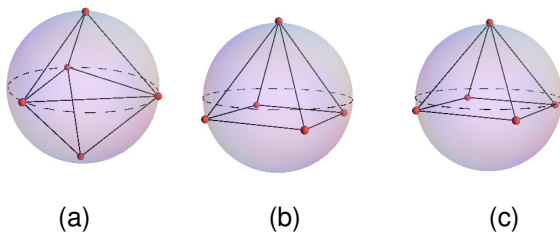
Optimal five point log and Riesz  $s$ -energy code on  $\mathbb{S}^2$ 

Figure: 'Optimal' 5-point codes on  $\mathbb{S}^2$ : (a) bipyramid BP, (b) optimal square-base pyramid SBP ( $s = 1$ ), (c) 'optimal' SBP ( $s = 16$ ).

- $s = 0$ : P. Dragnev, D. Legg, and D. Townsend, (2002) **(referred to by Ed Saff as "Mastodon" theorem)**;
- $s = -1$ : X. Hou, J. Shao, (2011), computer-aided proof;
- $s = 1, 2$ : R. E. Schwartz (2013), computer-aided proof;
- Bondarenko-Hardin-Saff (2014), As  $s \rightarrow \infty$ , any optimal  $s$ -energy codes of 5 limit is a square pyramid with base in the Equator;
- $0 < s < 15.04..$ : R. E. Schwartz (2018).

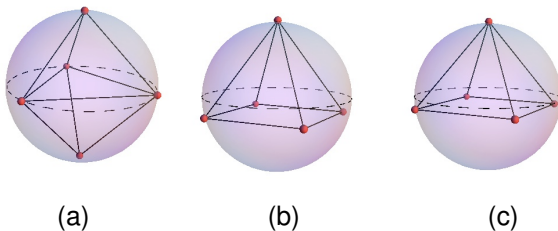
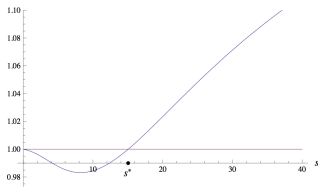
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Melnik et.al. 1977  $s^* = 15.04 \dots ?$

Figure: 5 points energy ratio

# “Mastodon” Theorem on $\mathbb{S}^3$ and $\mathbb{S}^4$ (Dragnev - 2016)

## Definition

Two vertices  $\mathbf{x}_i$  and  $\mathbf{x}_j$  are called *mirror related* (we write  $\mathbf{x}_i \sim \mathbf{x}_j$ ), if  $|\mathbf{x}_i - \mathbf{x}_k| = |\mathbf{x}_j - \mathbf{x}_k|$ , for every  $k \neq i, j$ .

## Theorem (Characterization of $(d + 3)$ Log-stationary configurations)

*A log-stationary configuration is either (a) degenerate; (b) there exists a vertex with all edges stemming out being equal; or (c) every vertex is mirror related to another vertex.*

## Remark

*Mirror relation is equivalence relation and an equivalence class forms a regular simplex in the spanning affine hyperspace.*

## Theorem (Dragnev - 2016)

*The  $(d + 3)$ -Log-optimal configuration in  $\mathbb{S}^1$ ,  $\mathbb{S}^2$ ,  $\mathbb{S}^3$ ,  $\mathbb{S}^4$ , is two orthogonal simplexes of type  $\{2, 2\}$ ,  $\{2, 3\}$ ,  $\{3, 3\}$ ,  $\{3, 4\}$  respectively.*

# “Mastodon” Theorem on $\mathbb{S}^{d-1}$ (Musin, D. - 2020)

## Theorem (Main Theorem 1)

*Up to orthogonal transform, every relative minimum of the logarithmic energy  $E_{\log}(X)$  of  $d + 2$  points on  $\mathbb{S}^{d-1}$  consists of two regular simplexes of cardinality  $m \geq n > 1$ ,  $m + n = d + 2$ , such that these simplexes are orthogonal to each other. The global minimum occurs when  $m = n$  if  $d$  is even and  $m = n + 1$  otherwise.*

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## Remark

*The only two other classes of minimal energy configurations are the regular simplex ( $d + 1$  points on  $\mathbb{S}^{d-1}$ ) and the regular cross polytope ( $2d$  points on  $\mathbb{S}^{d-1}$ ), both of which are universally optimal.*



Stationary Configurations of  $d + 2$  points on  $\mathbb{S}^{d-1}$ 

## Definition

A point configuration is called **degenerate** if it is contained in an affine hyperplane. (Ex. Pentagon on  $\mathbb{S}^2$ )

## Theorem (Non-degenerate, non-equidistant case)

*Let  $N = d + 2$  and  $X = \{x_1, \dots, x_N\}$  be a non-degenerate stationary logarithmic configuration on  $\mathbb{S}^{d-1}$ . Suppose there is no point  $x \in X$  that is equidistant to all other points in  $X$ . Then  $X$  can be split into two sets such that these sets are vertices of two regular orthogonal simplexes with the centers of mass in the center of  $\mathbb{S}^{d-1}$ .*

## Remark

*Strengthens 2016 Characterization theorem significantly.*

# Stationary Configurations of $d + 2$ points on $\mathbb{S}^{d-1}$

Given potential interaction function  $h : [-1, 1] \rightarrow \overline{\mathbb{R}}$   $h$ -energy is

$$E_h(X) := \sum_{1 \leq i \neq j \leq N} h(x_i \cdot x_j).$$

## Theorem (Degenerate Case - $h$ -energy)

*Let  $X$  be any degenerate configuration,  $N \geq d + 2$ , and  $h : [-1, 1] \rightarrow \mathbb{R}$  be a strictly convex potential function. Then there exists a continuous perturbation that decreases the  $h$ -energy  $E_h(X)$ .*

## Theorem (Equidistant case)

*A non-degenerate stationary log-energy configuration of type  $\{1, 1, \dots, k, l\}$ , where  $1 + 1 + \dots + k + l = d + 2$  is a saddle point. Moreover, there is a continuous perturbation that decreases the logarithmic energy of the  $\{1, k, l\}$  part of the configuration to either  $\{k + 1, l\}$  or  $\{k, l + 1\}$ . Sequence of such perturbations leads to relative minima as described in Main Theorem.*

## Auxiliary Results

Using Lagrange Multipliers method to *logarithmic energy*

$$E_{\text{Log}}(X) := -\frac{1}{2} \sum_{1 \leq i \neq j \leq N} \log(x_i \cdot x_i - 2x_i \cdot x_j + x_j \cdot x_j),$$

and differentiating yields

$$\sum_{j \neq i} \frac{x_i - x_j}{r_{i,j}} = \lambda_i x_i \quad i = 1, \dots, N, \text{ where } r_{ij} := 1 - x_i \cdot x_j.$$

Taking inner product of both sides with  $x_i$  one obtains  $\lambda_i = N - 1$ , or

$$\sum_{j \neq i} \frac{x_i - x_j}{r_{i,j}} = (N - 1) x_i, \quad i = 1, \dots, N. \quad (1)$$

Summing (1) implies that the centroid lies at the origin, and hence

$$\sum_j r_{ij} = N, \quad i = 1, \dots, N. \quad (2)$$

# Auxiliary Results - Rank Lemma

Let

$$B = (b_{ij}), \quad b_{ij} := \frac{1}{r_{ij}}, \quad b_{ii} := N - 1 - \sum_{j \neq i} b_{ij},$$

$$A = (a_{ij}), \quad \text{where } a_{ij} := c - b_{ij}, \quad c := \frac{N-1}{N}.$$

## Lemma

Let  $X = \{x_1, \dots, x_N\}$  be a stationary logarithmic configuration on  $\mathbb{S}^{d-1}$  that is non-degenerate. Then

$$\text{rank}(A) \leq N - d - 1, \quad \sum_{j=1}^N a_{ij} = 0, \quad i = 1, \dots, N.$$

If  $N = d + 2$ , then  $\text{rank}(A) = 1$ .

# Proof of the Rank Lemma

Let  $X := [x_1, \dots, x_N]^T$ . The force equations (1) and (2) imply that

$$\sum_{j=1}^N b_{ij} x_j = 0, \quad \sum_{j=1}^N b_{ij} = N - 1.$$

In other words,  $BX = 0$  and  $B\mathbf{1} = (N - 1)\mathbf{1}$ , where  $\mathbf{1}$  denotes the  $N$ -dimensional column-vector of ones.

As  $X$  is non-degenerate, we have  $\text{rank } X = d$ . Therefore, the column-vectors of  $X$  are linearly independent.

Since  $\mathbf{1}$  is eigenvector of  $B$  with an eigenvalue of  $N - 1$ , it is linearly independent to the columns of  $X$  (eigenvectors with eigenvalue 0).

The lemma follows from the rank-nullity theorem applied to  $A[X, \mathbf{1}] = 0$ .



# Auxiliary Results - $N = d + 2$

The following lemma elaborates on the case when  $N = d + 2$ .

## Lemma

Let  $N = d + 2$  and  $X = \{x_1, \dots, x_N\}$  be a non-degenerate stationary logarithmic configuration on  $\mathbb{S}^{d-1}$ . Without loss of generality we may assume that  $a_{1i} \geq 0$  for  $i = 1, \dots, k$  and  $a_{1i} < 0$  for  $i = k + 1, \dots, N$ . Let

$$a_i = \sqrt{a_{ii}}, \quad i = 1, \dots, k; \quad a_i = -\sqrt{a_{ii}}, \quad i = k + 1, \dots, N.$$

Then

$$a_{ij} = a_i a_j, \quad a_1 + \dots + a_N = 0,$$

$$c - a_i a_j \geq \frac{1}{2}, \quad \text{for all } i \neq j,$$

$$\sum_{j \neq i} \frac{1}{c - a_i a_j} = N, \quad i = 1, \dots, N. \quad (3)$$

# Auxiliary Results - Supplemental Theorem

If  $a_i = 0$  then the  $i$ -th row and  $i$ -th column in  $A$  are zero  $x_i$  is equidistant to all other points  $x_j$ . So,  $a_i \neq 0$  for all  $i = 1, \dots, N$ .

## Theorem (Supplemental)

Let  $a_1, \dots, a_N$  be real numbers that satisfy the following assumptions

$$a_1 \geq \dots \geq a_k > 0 > a_{k+1} \geq \dots \geq a_N, \quad a_1 + \dots + a_N = 0,$$

$$\sum_{j \neq i} \frac{1}{c - a_i a_j} = N, \quad i = 1, \dots, N, \quad c - a_i a_j > 0, \quad \text{for all } i \neq j,$$

where  $c := \frac{N-1}{N}$ . Then

$$a_1 = \dots = a_k, \quad a_{k+1} = \dots = a_N.$$

## Auxiliary Results - Technical Lemma

## Lemma (Technical)

Suppose  $a_1, \dots, a_N$  are as in Supplemental Theorem. Then for all  $i = 1, \dots, N$  we have

$$T_i := \sum_{j \neq i} \frac{c - a_j^2}{c - a_i a_j} = N - 2. \quad (4)$$

and

$$|a_i| < \sqrt{c}, \quad i = 1, \dots, N. \quad (5)$$



## Proof of the Technical Lemma - Eq. (4)

Denote

$$Q_i := \sum_{j \neq i} \frac{1}{c - a_i a_j}, \quad R_i := \sum_{j \neq i} \frac{a_j}{c - a_i a_j}, \quad S_i := \sum_{j \neq i} \frac{a_j^2 - a_j a_i}{c - a_i a_j}.$$

By the assumption  $Q_i = N$  for all  $i$ , we get (recall  $a_i \neq 0$ )

$$N - 1 = \sum_{j \neq i} \frac{c - a_i a_j}{c - a_i a_j} = c Q_i - a_i R_i = N - 1 - a_i R_i,$$

or  $R_i = 0$ . Along with  $a_i = -(a_1 + \dots + a_{i-1} + a_{i+1} + \dots + a_N)$

$$a_i = (c - a_i^2) \sum_{j \neq i} \frac{a_j}{c - a_i a_j} - \sum_{j \neq i} a_j = a_i \sum_{j \neq i} \frac{a_j^2 - a_j a_i}{c - a_i a_j} = a_i S_i,$$

or  $S_i = 1$  and subsequently

$$N - 2 = \sum_{j \neq i} \frac{c - a_i a_j}{c - a_i a_j} - S_i = \sum_{j \neq i} \frac{c - a_j^2}{c - a_i a_j} = T_i$$

## Proof of the Technical Lemma - Ineq. (5)

W.l.g.  $|a_1| \geq |a_i|$ . From  $S_i = 1$  we have

$$1 = \sum_{j \neq i} \frac{a_j^2 - a_i a_j}{c - a_i a_j} = \frac{a_1^2 - a_i a_1}{c - a_i a_1} + \sum_{2 \leq j \neq i} \frac{a_j^2 - a_i a_j}{c - a_i a_j}.$$

Then

$$\sum_{2 \leq j \neq i} \frac{a_j^2 - a_i a_j}{c - a_i a_j} = \frac{c - a_1^2}{c - a_i a_1}, \quad i = 2, \dots, N.$$

Therefore,

$$\sum_{i=2}^N \sum_{2 \leq j \neq i} \frac{a_j^2 - a_i a_j}{c - a_i a_j} = \sum_{i>j=2}^N \frac{(a_i - a_j)^2}{c - a_i a_j} = (c - a_1^2) \sum_{i=2}^N \frac{1}{c - a_i a_1} = (c - a_1^2) Q_1.$$

Since  $Q_1 = N$  and by the assumption  $c - a_i a_j > 0$ , we have

$$c - a_1^2 = \frac{1}{N} \sum_{i>j=2}^N \frac{(a_i - a_j)^2}{c - a_i a_j} > 0. \quad (6)$$

Thus, (6) implies  $c - a_1^2 > 0$ . □

## Proof of Supplemental Theorem

Let

$$F(t) := \sum_{j=1}^N \frac{c - a_j^2}{c - ta_j}.$$

Then Technical Lemma implies that for all  $i = 1, \dots, N$

$$F(a_i) = N - 1. \quad (7)$$

Since

$$F''(t) = 2 \sum_j \frac{(c - a_j^2) a_j^2}{(c - ta_j)^3},$$

by Technical Lemma again we have  $F''(t) > 0$  for  $t \in (-\sqrt{c}, \sqrt{c})$ . Hence  $F(t)$  is a convex function in this interval. Therefore, the equation  $F(t) = N - 1$  has at most two solutions. By assumptions we have  $a_i > 0$  for  $i = 1, \dots, k$  and  $a_i < 0$ , for  $i = k + 1, \dots, N$ . Thus, (7) yields that all positive  $a_i$  are equal and all negative  $a_i$  are equal too.  $\square$

# Proofs of Degenerate, Equidistant, and Relative Minima cases

Even more complex and involved :-)

# Degenerate case - $h$ -energy

## Theorem (Degenerate Case)

Let  $X$  be a degenerate configuration,  $N \geq d + 2$ , and  $h : [-1, 1] \rightarrow \mathbb{R}$  be a strictly convex potential function. Then there exists a continuous perturbation that decreases the  $h$ -energy  $E_h(X)$ .

## Proof.

$$x_1 = (r, \sqrt{1 - r^2}, 0, \dots, 0), x_2 = (r, -\sqrt{1 - r^2}, 0, \dots, 0)$$

$$x_j = (c_{j1}, c_{j2}, c_{j3}, \dots, 0), j = 3, \dots, N,$$

where  $c_{32} \neq 0$ . Perturb to  $\tilde{X}$

$$\tilde{x}_1 = (r, \sqrt{1 - r^2} \cos \theta, 0, \dots, \sqrt{1 - r^2} \sin \theta),$$

$$\tilde{x}_2 = (r, -\sqrt{1 - r^2} \cos \theta, 0, \dots, -\sqrt{1 - r^2} \sin \theta).$$

Then  $E_h(X) > E_h(\tilde{X})$



# Equidistant case

## Theorem (Equidistant case)

*A non-degenerate stationary log-energy configuration of type  $\{1, 1, \dots, k, \ell\}$ , where  $1 + 1 + \dots + k + \ell = d + 2$  is a saddle point. Moreover, there is a continuous perturbation that decreases the logarithmic energy of the  $\{1, k, \ell\}$  part of the configuration to either  $\{k + 1, \ell\}$  or  $\{k, \ell + 1\}$ . Sequence of such perturbations leads to relative minima as described in Main Theorem.*

## Equidistant case proof

## Proof.

Let  $X = \{1, k, \ell\}$  with  $x_N \cdot x_i = -1/(N-1)$ . Denote  $x_i = (y_i, \frac{-1}{N-1})$ ,  $z_i := (N-1)y_i/\sqrt{N(N-2)}$ ,  $z_i \in \mathbb{S}^{d-2}$  satisfies force equation.

$$Y := \{(\sqrt{1 - 1/(k+m)^2} y_i, 0_{m-1}, -1/(k+m))\},$$

$$Z := \{(0_{k-1}, \sqrt{1 - 1/(k+m)^2} z_j, -1/(k+m))\}$$

Perturb to

$$\tilde{Y}_t = \left\{ \left( \sqrt{1 - (mt + 1/(k+m))^2} y_i, 0_{m-1}, -1/(k+m) - mt \right) \right\}_{i=1}^k$$

$$\tilde{Z}_t = \left\{ \left( 0_{k-1}, \sqrt{1 - (kt - 1/(k+m))^2} z_j, -1/(k+m) + kt \right) \right\}_{j=1}^m.$$

Then  $E_h(\tilde{X}_t)$  has local max at  $t = 0$  and decreases to  $\{k, \ell + 1\}$  or  $\{k + 1, \ell\}$ . □

## Relative minima case

### Theorem (Equidistant case)

*Let  $X = \{k, \ell\}$  a configuration of two orthogonal simplexes  $X_k$  and  $X_\ell$ . Any perturbation will increase the energy locally.*

We need two inequalities.



## Relative minima case - Inequality 1

## Lemma (H1)

Let  $A = (a_{ij})$  be an  $m \times m$  matrix,  $m \geq 3$ , such that

(a)  $a_{ii} = 0$ ,  $i = 1, \dots, m$ ;

(b)  $\sum_{j=1}^m a_{ij} = 0$ .

Then the following inequality holds

$$\sum_{1 \leq i < j \leq m} (a_{ij} + a_{ji})^2 \geq \frac{1}{m-2} \sum_{j=1}^m x_j^2, \quad \text{where } x_j := \sum_{i=1}^m a_{ij}. \quad (8)$$

# Proof of Inequality 1: part 1

For all  $i, j = 1, \dots, m$  define

$$\beta_{ij} := \frac{1}{m^2 - 2m} x_i + \frac{m-1}{m^2 - 2m} x_j, \quad i \neq j, \quad \text{and} \quad \beta_{ii} = 0.$$

Since  $\sum_{j=1}^m x_j = 0$ , we have  $\sum_{j=1}^m \beta_{ij} = 0$  and  $\sum_{i=1}^m \beta_{ij} = x_j$ , i.e.

$$\sum_{j=1}^m \beta_{ij} = \sum_{j=1}^m a_{ij} \quad \text{and} \quad \sum_{i=1}^m \beta_{ij} = \sum_{i=1}^m a_{ij}.$$

Let  $\tilde{a}_{ij} := a_{ij} - \beta_{ij}$ . Then

$$\sum_i \tilde{a}_{ij} = \sum_j \tilde{a}_{ij} = 0.$$

# Proof of Inequality 1: part 2

Consider  $t_{ij} := a_{ij} + a_{ji} = w_{ij} + \beta_{ij} + \beta_{ji}$ , where  $w_{ij} = \tilde{a}_{ij} + \tilde{a}_{ji}$ . Then  $t_{ij} = w_{ij} + \frac{x_i}{m-2} + \frac{x_j}{m-2}$ ,  $i \neq j$ , where  $\sum_i w_{ij} = \sum_j w_{ij} = 0$  (observe that  $t_{ii} = 0$ ). Then

$$\sum_{i < j} t_{ij}^2 = \sum_{i < j} \left( w_{ij} + \frac{x_i}{m-2} + \frac{x_j}{m-2} \right)^2 = \sum_{i < j} w_{ij}^2 + \frac{1}{m-2} \sum_{i=1}^m x_i^2,$$

which implies (8).

## Relative minima case - Inequality 2

## Lemma (H2)

Given an  $m \times n$  matrix  $F = (f_{ij})$  and an  $n \times m$  matrix  $G = (g_{ij})$  such that  $\sum_{j=1}^n f_{ij} = 0$  for all  $i = 1, \dots, m$  and  $\sum_{j=1}^m g_{ij} = 0$  for all  $i = 1, \dots, n$ . Then we have

$$\sum_{i=1}^n \sum_{j=1}^m (f_{ij} + g_{ji})^2 \geq \frac{1}{m} \sum_{j=1}^n y_j^2 + \frac{1}{n} \sum_{i=1}^m z_i^2,$$

$$y_j := \sum_{i=1}^m f_{ij}, \quad z_i := \sum_{j=1}^n g_{ji}.$$

## Proof of Inequality 2

Let

$$\tilde{f}_{ij} := f_{ij} - \frac{y_j}{m} \quad \text{and} \quad \tilde{g}_{ij} := g_{ij} - \frac{z_i}{n}.$$

Since  $\sum_j y_j = \sum_i z_i = 0$ , we have  $\sum_{i,j} (\tilde{f}_{ij} + \tilde{g}_{ji}) = 0$ . Let  $t_{ij} := \tilde{f}_{ij} + \tilde{g}_{ji}$ . Observe that

$$\sum_{i=1}^m t_{ij} = \sum_{j=1}^n t_{ij} = 0.$$

From

$$f_{ij} + g_{ji} = \frac{y_j}{m} + \frac{z_i}{n} + t_{ij}.$$

one derives that

$$\sum_{i=1}^m \sum_{j=1}^n (f_{ij} + g_{ji})^2 = \sum_{i=1}^m \sum_{j=1}^n \left( \frac{y_j}{m} + \frac{z_i}{n} + t_{ij} \right)^2 = \sum_{i=1}^m \sum_{j=1}^n t_{ij}^2 + \frac{1}{m} \sum_{j=1}^n y_j^2 + \frac{1}{n} \sum_{i=1}^m z_i^2,$$

which completes the proof.

## Relative minima case - Proof (part 1)

Perturb the two orthogonal simplexes ( $k + \ell = d + 2$ )

$$X_k = \{x_1, x_2, \dots, x_k\}, \quad X_\ell = \{x_{k+1}, x_{m+2}, \dots, x_{k+\ell}\},$$

to  $Y = Y_k \cup Y_\ell$ , where  $y_i := x_i + h_i$ ,  $\|h_i\| < \epsilon$ . Since  $\|x_i\| = \|y_i\| = 1$ , we have  $2x_i \cdot h_i = -\|h_i\|^2$ ,  $1 - y_i \cdot y_j = (1 - x_i \cdot x_j)(1 - z_{i,j})$ , where

$$z_{i,j} := \begin{cases} \frac{k-1}{k}(x_i \cdot h_j + x_j \cdot h_i + h_i \cdot h_j), & 1 \leq i \neq j \leq k \\ x_i \cdot h_j + x_j \cdot h_i + h_i \cdot h_j, & i \leq k < j \text{ or } j \leq k < i \\ \frac{\ell-1}{\ell}(x_i \cdot h_j + x_j \cdot h_i + h_i \cdot h_j), & k < i \neq j \leq k + \ell. \end{cases} \quad (9)$$

Clearly  $|z_{i,j}| < 2\epsilon + O(\epsilon^2)$ . We find

$$2[E_{\log}(Y) - E_{\log}(X)] = \sum_{1 \leq i \neq j \leq k+\ell} \left( z_{i,j} + \frac{z_{i,j}^2}{2} \right) + O(\epsilon^3). \quad (10)$$

## Relative minima case - Proof (part 2)

W.l.g.  $x_i = (p_i, 0)$ ,  $h_i = (a_i, b_i)$ ,  $x_{k+j} = (0, q_j)$ ,  $h_{k+j} = (c_j, d_j)$ , where  $p_i, a_i, c_j \in \mathbb{R}^{k-1}$  and  $q_j, b_i, d_j \in \mathbb{R}^{\ell-1}$ . Straight-forward calculations show that the linear in  $\epsilon$  term in (10) vanishes. The quadratic term is

$$\begin{aligned}
 D &:= \sum_{1 \leq i \neq j \leq k+\ell} \frac{h_i \cdot h_j}{1 - x_i \cdot x_j} + \frac{1}{2} \sum_{1 \leq i \neq j \leq k+\ell} \left( \frac{x_i \cdot h_j + x_j \cdot h_i}{1 - x_i \cdot x_j} \right)^2 \\
 &= \left\| \sum_{i=1}^{k+\ell} h_i \right\|^2 - \frac{1}{k} \left( \left\| \sum_{i=1}^k a_i \right\|^2 + \left\| \sum_{i=1}^k b_i \right\|^2 \right) - \frac{1}{\ell} \left( \left\| \sum_{j=1}^{\ell} c_j \right\|^2 + \left\| \sum_{j=1}^{\ell} d_j \right\|^2 \right) \\
 &+ \left( \frac{k-1}{k} \right)^2 \sum_{1 \leq i < j \leq k} (p_i \cdot a_j + p_j \cdot a_i)^2 + \left( \frac{\ell-1}{\ell} \right)^2 \sum_{1 \leq i < j \leq \ell} (q_i \cdot d_j + q_j \cdot d_i)^2 \\
 &+ \sum_{i=1}^k \sum_{j=1}^{\ell} (p_i \cdot c_j + q_j \cdot b_i)^2.
 \end{aligned} \tag{11}$$

## Relative minima case - Proof (part 3)

By extracting another  $O(\epsilon^3)$  term we may reduce the condition  $2x_i \cdot h_i = -\|h_i\|^2$  to  $x_i \cdot h_i = 0$ .

Thus, in this case we shall reduce the theorem to proving the inequalities

$$D_1 := \left(\frac{k-1}{k}\right)^2 \sum_{1 \leq i < j \leq k} (p_i \cdot a_j + p_j \cdot a_i)^2 - \frac{1}{k} \left\| \sum_{i=1}^k a_i \right\|^2 \geq 0 \quad (12)$$

$$D_2 := \left(\frac{\ell-1}{\ell}\right)^2 \sum_{1 \leq i < j \leq \ell} (q_i \cdot d_j + q_j \cdot d_i)^2 - \frac{1}{\ell} \left\| \sum_{j=1}^{\ell} d_j \right\|^2 \geq 0 \quad (13)$$

and

$$D_3 := \sum_{i=1}^k \sum_{j=1}^{\ell} (p_i \cdot c_j + q_j \cdot b_i)^2 - \frac{1}{k} \left\| \sum_{i=1}^k b_i \right\|^2 - \frac{1}{\ell} \left\| \sum_{j=1}^{\ell} c_j \right\|^2 \geq 0 \quad (14)$$

provided  $\{p_1, \dots, p_k\}$  and  $\{q_1, \dots, q_{\ell}\}$  are orthogonal  $k$ - and  $\ell$ -simplexes and  $p_i \cdot a_i = 0$  and  $q_j \cdot d_j = 0$ .



## Relative minima case - Proof (part 4)

Embed the first simplex  $X_k = \{p_1, \dots, p_k\}$  in the hyperplane of  $\mathbb{R}^k$  that is orthogonal to  $(1, 1, \dots, 1)$ . Similarly, we embed the second simplex  $X_\ell = \{q_1, \dots, q_\ell\}$  in  $\mathbb{R}^\ell$ . Thus, we embed  $X_k \cup X_\ell \subset \mathbb{R}^k \times \mathbb{R}^\ell$ .

Let  $w_k = (\frac{1}{k}, \frac{1}{k}, \dots, \frac{1}{k}) \in \mathbb{R}^k$ ,  $\tilde{p}_i := e_i - w_k$ , so  $p_i = \sqrt{\frac{k}{k-1}} \tilde{p}_i$ . Similarly, if  $\tilde{q}_j := e_j - w_\ell$ , then  $q_j = \sqrt{\frac{\ell}{\ell-1}} \tilde{q}_j$ .

For the perturbation vectors  $a_i, b_i, c_j, d_j$ , we have

$$\sum_{j=1}^k a_{ij} = 0, \quad \sum_{j=1}^{\ell} b_{ij} = 0, \quad \sum_{j=1}^k c_{ij} = 0, \quad \sum_{j=1}^{\ell} d_{ij} = 0.$$

$p_i \cdot a_i = 0 \Rightarrow a_{ii} = 0$ ,  $q_j \cdot d_j = 0 \Rightarrow a_{jj} = 0$ . Using  $\tilde{p}_i \cdot a_j = a_{ji}$  (12) is

$$\sum_{1 \leq i < j \leq k} (a_{ij} + a_{ji})^2 \geq \frac{1}{k-1} \sum_{j=1}^k \left( \sum_{i=1}^k a_{ij} \right)^2,$$

which follows from Lemma (H1). Similarly one gets (13).

## Relative minima case - Proof (part 5)

Equality in (12) and (13) holds iff  $a_{ij} + a_{ji} = 0$  and  $d_{ij} + d_{ji} = 0$  respectively, which is equivalent to

$$p_i \cdot a_j + p_j \cdot a_i = 0, q_j \cdot d_i + q_i \cdot d_j = 0, \sum a_i = 0, \sum d_j = 0.$$

Lemma (H2) is used to derive the inequality (14). We have that

$$p_i \cdot c_j + q_j \cdot b_i = \sqrt{\frac{k}{k-1}} c_{ji} + \sqrt{\frac{\ell}{\ell-1}} b_{ij}$$

with the substitution  $f_{ij} = \sqrt{\frac{\ell}{\ell-1}} b_{ij}$  and  $g_{ji} = \sqrt{\frac{k}{k-1}} c_{ji}$  we re-write (14) as

$$\sum_{i=1}^{\ell} \sum_{j=1}^k (f_{ij} + g_{ij})^2 \geq \frac{1}{k} \frac{\ell-1}{\ell} \sum_{j=1}^{\ell} y_j^2 + \frac{1}{\ell} \frac{k-1}{k} \sum_{i=1}^k z_i^2,$$

which clearly follows from Lemma (H2). Moreover, equality occurs if and only if  $p_i \cdot c_j + q_j \cdot b_i = 0$ ,  $\sum c_i = 0$ , and  $\sum b_j = 0$ .

## Relative minima case - Proof (part 6)

So, the quadratic in  $\epsilon$  term  $D > 0$ , or  $E_{\log}(Y) - E_{\log}(X) > 0$ , for any perturbation vectors  $\{a_i, b_i, c_i, d_i\}$  ( $p_i \cdot a_i = 0$ ,  $q_j \cdot d_j = 0$ ), except when

$$p_i \cdot a_j + p_j \cdot a_i = 0, \quad q_j \cdot d_i + q_i \cdot d_j = 0, \quad p_i \cdot c_j + q_j \cdot b_i = 0,$$

$$\sum_{i=1}^k a_i = 0, \quad \sum_{i=1}^k b_i = 0, \quad \sum_{j=1}^{\ell} c_j = 0, \quad \sum_{j=1}^{\ell} d_j = 0.$$

Utilizing (9) and these conditions, one simplifies (10) to

$$\begin{aligned} 2[E_{\log}(Y) - E_{\log}(X)] &= \frac{(k-1)^2}{2k^2} \sum_{1 \leq i \neq j \leq k} (a_i \cdot a_j)^2 + \frac{(\ell-1)^2}{2\ell^2} \sum_{1 \leq i \neq j \leq \ell} (d_i \cdot d_j)^2 \\ &\quad + \sum_{i=1}^k \sum_{j=1}^{\ell} (b_i \cdot c_j)^2 + O(\epsilon^5). \end{aligned}$$

Clearly, the quartic term will be positive, unless all inner products vanish, in which case we easily derive that  $a_i = c_j = 0$  and  $b_i = d_j = 0$  for all  $i = 1, \dots, k$  and  $j = 1, \dots, \ell$ . This completes the proof.  $\square$

**THANK YOU!**